## States on Hilbert Lattices (Vienna, April 22, 2008)

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## Early Ideas and Results

Ancient Result: There is an isomorphism between a Hilbert lattice (a complete atomic orthomodular lattice which satisfies the superpostion princile and has > 2 atoms) and the set of all closed subspaces of a Hilbert space.

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If we wanted to connect a Hilbert lattice with measured physical states of a quantum system described by a Hilbert space equation, we should impose quantum states on the lattice.
What is a quantum state on a lattice? Are there classical states?

## States

Definition. A state on a lattice $L$ is a function $m: \mathrm{L} \longrightarrow[0,1]$ (for real interval $[0,1]$ ) such that $m(1)=1$ and $a \perp b \Rightarrow m(a \cup b)=m(a)+m(b)$, where $a \perp b$ means $a \leq b^{\prime}$.

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Definiton. A nonempty set $S$ of states on L is called a strong set of classical states if
$(\exists m \in S)(\forall a, b \in \mathrm{~L})((m(a)=1 \Rightarrow m(b)=1) \Rightarrow a \leq b)$
and a strong set of quantum states if
$(\forall a, b \in \mathrm{~L})(\exists m \in S)((m(a)=1 \Rightarrow m(b)=1) \Rightarrow a \leq b)$

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Some of these states might be useful.

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■ full but not strong — they just show that there are othomodular lattices that are not Hilbert ones - with our algorithms and programs we can always generate a pile of them if some application pops up

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■ Frederic F. Shultz, J. Comb. Theory A 17, 317 (1974) $\longrightarrow$ Mirko Navara, Int. J. Theor. Phys. 47, 36 (2008)

## Efficiency

A good scenario: states generate Hilbert lattice equations - introduced by Radosław Godowski in 1981

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However, they might be simplifed when considered as Greechie diagrams and hypergraphs

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We call the diagrams MMP diagrams when they do not refer to any structure-when they are just dots and lines; vertices and edges; atoms and blocks.

## MMP definitions

They are graphs defined as follows:
(1) Every atom (vertex, point) belongs to at least one block (edge, line).
(2) If there are at least two atoms then every block is at least 2-element.
(3) Every block which intersects with another block is at least 3-element.
(4) Every pair of different blocks intersects in at most one (two, three) atom(s).
(5) Smallest loops are of order $1(2,3,4,5)$

## Finding States

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If $m$ is a state, then each 3 -atom block with atoms $a, b, c$ imposes the following constraints:

$$
\begin{aligned}
m(a)+m(b)+m(c) & =1 \\
m\left(a^{\prime}\right)+m(a) & =1 \\
m\left(b^{\prime}\right)+m(b) & =1 \\
m\left(c^{\prime}\right)+m(c) & =1 \\
m(x) & \geq 0, \quad x=a, b, c, a^{\prime}, b^{\prime}, c^{\prime}
\end{aligned}
$$

## Lattice Equation

A condensed state equation is an abbreviated version of a lattice equation constructed as follows: all (orthogonality) hypotheses are discarded, all meet symbols, $\cap$, are changed to + , and all join symbols, $\cup$, are changed to juxtaposition.

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\begin{aligned}
& \text { E.g. } a \perp d \perp b \perp e \perp c \perp f \perp a \Rightarrow \\
& (a \cup d) \cap(b \cup e) \cap(c \cup f)=(d \cup b) \cap(e \cup c) \cap(f \cup a)
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which, in turn, can be expressed by the condensed state equation
$a d+b e+c f=d b+e c+f a$.

## States $\rightarrow$ Equations

When our program finds the states then we obtain various constraints, such as:

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m(B)+m(C)+m(1) \leq 1 ; m(2)+m(E)+m(8)=1
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Replacing the atoms with variables, the final condensed state equation becomes:
$a b+c d+e f+g h=b g+f c+a d+h e$

