

Numeričke metode

Aleksandar Maksimović
IRB

Sustavi linearnih jednačbi

- ♦ Sustav n jednadžbi i n nepoznanica

$$a_{11}x_1 + a_{12}x_2 + \cdots + a_{1j}x_j + \cdots + a_{1n}x_n = b_1$$

$$a_{21}x_1 + a_{22}x_2 + \cdots + a_{2j}x_j + \cdots + a_{2n}x_n = b_2$$

$$\vdots \qquad \vdots \qquad \vdots \qquad \vdots \qquad \vdots$$

$$a_{i1}x_1 + a_{i2}x_2 + \cdots + a_{ij}x_j + \cdots + a_{in}x_n = b_i$$

$$\vdots \qquad \vdots \qquad \vdots \qquad \vdots \qquad \vdots$$

$$a_{n1}x_1 + a_{n2}x_2 + \cdots + a_{nj}x_j + \cdots + a_{nn}x_n = b_n$$

Matrica $A = [a_{ij}]_{i,j=1}^n \in \mathbb{R}^{n \times n}$ je **matrica sustava**, a njeni elementi su **koeficijenti sustava**. Vektor $b = [b_i]_{i=1}^n \in \mathbb{R}^n$ je **vektor desne strane** sustava. Treba odrediti **vektor nepoznanica** $x = [x_i]_{i=1}^n \in \mathbb{R}^n$ tako da vrijedi $Ax = b$.

Primjer 1

Interpolacijski polinom

$$p(x) = a_0 + a_1x + a_2x^2 + \cdots + a_nx^n = \sum_{j=0}^n a_jx^j$$

nepoznati su koeficijenti, određujemo ih poznatih vrijednosti y

$$p(x_i) = y_i, \quad i = 0, \dots, n,$$

$$a_0 + a_1x_0 + a_2x_0^2 + \cdots + a_{n-1}x_0^{n-1} + a_nx_0^n = y_0$$

$$a_0 + a_1x_1 + a_2x_1^2 + \cdots + a_{n-1}x_1^{n-1} + a_nx_1^n = y_1$$

$$\vdots \quad \vdots \quad \vdots \quad \dots \quad \vdots \quad \vdots \quad \vdots$$

$$a_0 + a_1x_i + a_2x_i^2 + \cdots + a_{n-1}x_i^{n-1} + a_nx_i^n = y_i$$

$$a_0 + a_1x_n + a_2x_n^2 + \cdots + a_{n-1}x_n^{n-1} + a_nx_n^n = y_n.$$

Primjer 1

$$Va = y$$

$$\underbrace{\begin{bmatrix} 1 & x_0 & x_0^2 & x_0^3 & \cdots & x_0^{n-1} & x_0^n \\ 1 & x_1 & x_1^2 & x_1^3 & \cdots & x_1^{n-1} & x_1^n \\ \cdots & \cdots & \cdots & \cdots & \cdots & \cdots & \cdots \\ 1 & x_i & x_i^2 & x_i^3 & \cdots & x_i^{n-1} & x_i^n \\ \cdots & \cdots & \cdots & \cdots & \cdots & \cdots & \cdots \\ 1 & x_n & x_n^2 & x_n^3 & \cdots & x_n^{n-1} & x_n^n \end{bmatrix}}_V \underbrace{\begin{bmatrix} a_0 \\ a_1 \\ a_2 \\ \vdots \\ a_{n-1} \\ a_n \end{bmatrix}}_a = \underbrace{\begin{bmatrix} y_0 \\ y_1 \\ y_2 \\ \vdots \\ y_{n-1} \\ y_n \end{bmatrix}}_y.$$

Primjer 2

Promotrimo sljedeći rubni problem:

$$\begin{aligned} -\frac{d^2}{dx^2}u(x) &= f(x), \quad 0 < x < 1, \\ u(0) &= u(1) = 0. \end{aligned}$$

$$h = \frac{1}{n+1}, \quad x_i = ih, \quad i = 0, \dots, n+1.$$

$$u_i = u(x_i), \quad i = 0, \dots, n+1.$$

$$u_0 = u_{n+1} = 0$$

Primjer 2

$$u(x_i + h) = u(x_i) + u'(x_i)h + \frac{u''(x_i)}{2}h^2 + \frac{u'''(x_i)}{6}h^3 + \frac{u^{(4)}(x_i + \alpha_i)}{24}h^4$$

$$u(x_i - h) = u(x_i) - u'(x_i)h + \frac{u''(x_i)}{2}h^2 - \frac{u'''(x_i)}{6}h^3 + \frac{u^{(4)}(x_i + \zeta_i)}{24}h^4,$$

Kombinacijom ova dva Taylorova reda dobivamo drugu derivaciju

$$u_{i+1} + u_{i-1} = 2u_i + u''(x_i)h^2 + (u^{(4)}(x_i + \alpha_i) + u^{(4)}(x_i + \zeta_i))\frac{h^4}{24},$$

Druga derivacija metodom konačnih razlika

$$-u''(x_i) = \frac{-u_{i-1} + 2u_i - u_{i+1}}{h^2}$$

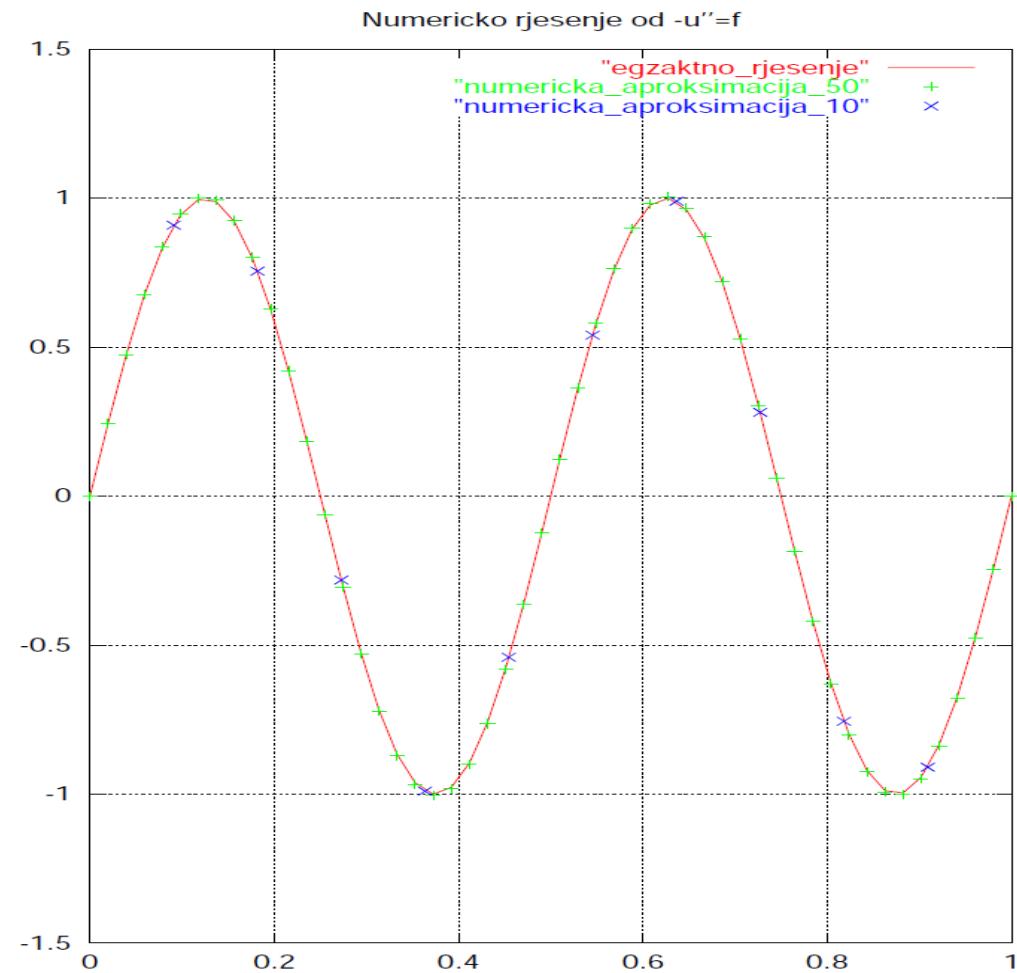
Primjer 2

$$\underbrace{\begin{bmatrix} 2 & -1 & & & \\ -1 & 2 & -1 & & \\ & -1 & 2 & -1 & \\ & & \cdots & \cdots & \cdots \\ & & & -1 & 2 & -1 \\ & & & & -1 & 2 & -1 \\ & & & & & -1 & 2 \end{bmatrix}}_{T_n} \underbrace{\begin{bmatrix} u_1 \\ u_2 \\ u_3 \\ \vdots \\ u_{n-2} \\ u_{n-1} \\ u_n \end{bmatrix}}_u = h^2 \underbrace{\begin{bmatrix} f_1 \\ f_2 \\ f_3 \\ \vdots \\ f_{n-2} \\ f_{n-1} \\ f_n \end{bmatrix}}_f.$$

Primjer 2

$$f(x) = 16\pi^2 \sin 4\pi x$$

$$n = 10 \text{ i } n = 50$$



Linearni sustav 2x2

$$\begin{aligned} 2x_1 - x_2 &= 1 \\ -x_1 + 2x_2 &= 1 \end{aligned}$$

$$x_1 = \frac{1}{2}(1 + x_2)$$

$$-\underbrace{\frac{1}{2}(1 + x_2)}_{x_1} + 2x_2 = 1, \quad \text{tj. } \frac{3}{2}x_2 = \frac{3}{2}, \quad \text{tj. } x_2 = 1, \quad x_1 = 1$$

Metoda supstitucije

$$\begin{array}{lcl} 5x_1 + x_2 + 4x_3 = 19 \\ 10x_1 + 4x_2 + 7x_3 = 39 \\ -15x_1 + 5x_2 - 9x_3 = -32 \end{array} \equiv \underbrace{\begin{bmatrix} 5 & 1 & 4 \\ 10 & 4 & 7 \\ -15 & 5 & -9 \end{bmatrix}}_{A = [a_{ij}]_{i,j=1}^3} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \end{bmatrix} = \underbrace{\begin{bmatrix} 19 \\ 39 \\ -32 \end{bmatrix}}_{b = [b_i]_{i=1}^3}.$$

Koristimo metodu supstitucija, odnosno eliminacija. Prvo iz prve jednadžbe izrazimo x_1 pomoću x_2 i x_3 , te to uvrstimo u zadnje dvije jednadžbe, koje postaju dvije jednadžbe s dvije nepoznanice (x_2 i x_3). Dobivamo

$$x_1 = \frac{1}{5} (19 - x_2 - 4x_3),$$

Metoda supstitucije

pa druga jednadžba sada glasi

$$\frac{10}{5} (19 - x_2 - 4x_3) + 4x_2 + 7x_3 = 39,$$

tj.

$$-\frac{10}{5} (x_2 + 4x_3) + 4x_2 + 7x_3 = 39 + \left(-\frac{10}{5} 19\right)$$

Dakle, efekt ove transformacije je ekvivalentno prikazan kao rezultat množenja prve jednadžbe s

$$-\frac{a_{21}}{a_{11}} = -\frac{10}{5} = -2$$

i zatim njenim dodavanjem (pribrajanjem) drugoj jednadžbi. Druga jednadžba sada glasi

$$2x_2 - x_3 = 1.$$

Metoda supstitucije

$$\underbrace{\begin{bmatrix} 5 & 1 & 4 \\ 10 & 4 & 7 \\ -15 & 5 & -9 \end{bmatrix}}_A \mapsto \underbrace{\begin{bmatrix} 1 & 0 & 0 \\ -2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}}_{L^{(2,1)}} \underbrace{\begin{bmatrix} 5 & 1 & 4 \\ 10 & 4 & 7 \\ -15 & 5 & -9 \end{bmatrix}}_A = \underbrace{\begin{bmatrix} 5 & 1 & 4 \\ 0 & 2 & -1 \\ -15 & 5 & -9 \end{bmatrix}}_{A^{(1)}} \quad A^{(1)} = \left[a_{ij}^{(1)} \right]_{i,j=1}^3$$

Nepoznanicu x_1 eliminiramo iz zadnje jednadžbe ako prvu pomnožimo s

$$\underbrace{\begin{bmatrix} 5 & 1 & 4 \\ 0 & 2 & -1 \\ -15 & 5 & -9 \end{bmatrix}}_{A^{(1)}} \mapsto \underbrace{\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 3 & 0 & 1 \end{bmatrix}}_{L^{(3,1)}} \underbrace{\begin{bmatrix} 5 & 1 & 4 \\ 0 & 2 & -1 \\ -15 & 5 & -9 \end{bmatrix}}_{A^{(1)}} = \underbrace{\begin{bmatrix} 5 & 1 & 4 \\ 0 & 2 & -1 \\ 0 & 8 & 3 \end{bmatrix}}_{A^{(2)}} \quad A^{(2)} = \left[a_{ij}^{(2)} \right]_{i,j=1}^3$$

Vektor desne strane je u ove dvije transformacije promijenjen u

Metoda supstitucije

Vektor desne strane je u ove dvije transformacije promijenjen u

$$\underbrace{\begin{bmatrix} 19 \\ 39 \\ -32 \end{bmatrix}}_b \mapsto \underbrace{\begin{bmatrix} 1 & 0 & 0 \\ -2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}}_{L^{(2,1)}} \begin{bmatrix} 19 \\ 39 \\ -32 \end{bmatrix} = \underbrace{\begin{bmatrix} 19 \\ 1 \\ -32 \end{bmatrix}}_{b^{(1)}}$$

$$\mapsto \underbrace{\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 3 & 0 & 1 \end{bmatrix}}_{L^{(3,1)}} \begin{bmatrix} 19 \\ 1 \\ -32 \end{bmatrix} = \underbrace{\begin{bmatrix} 19 \\ 1 \\ 25 \end{bmatrix}}_{b^{(2)}}.$$

Novi, ekvivalentni, sustav je $A^{(2)}x = b^{(2)}$, tj.

Metoda supstitucije

Novi, ekvivalentni, sustav je $A^{(2)}x = b^{(2)}$, tj.

$$5x_1 + x_2 + 4x_3 = 19$$

$$2x_2 - x_3 = 1$$

$$8x_2 - 3x_3 = 25,$$

$$-\frac{a_{32}^{(2)}}{a_{22}^{(2)}} = -4$$

$$5x_1 + x_2 + 4x_3 = 19$$

$$2x_2 - x_3 = 1$$

$$7x_3 = 21.$$

Metoda supstitucije

$$\underbrace{\begin{bmatrix} 5 & 1 & 4 \\ 0 & 2 & -1 \\ 0 & 8 & 3 \end{bmatrix}}_{A^{(2)}} \mapsto \underbrace{\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & -4 & 1 \end{bmatrix}}_{L^{(3,2)}} \underbrace{\begin{bmatrix} 5 & 1 & 4 \\ 0 & 2 & -1 \\ 0 & 8 & 3 \end{bmatrix}}_{A^{(2)}} = \underbrace{\begin{bmatrix} 5 & 1 & 4 \\ 0 & 2 & -1 \\ 0 & 0 & 7 \end{bmatrix}}_{A^{(3)}}$$

transformaciju vektora desne strane

$$\underbrace{\begin{bmatrix} 19 \\ 1 \\ 25 \end{bmatrix}}_{b^{(2)}} \mapsto \underbrace{\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & -4 & 1 \end{bmatrix}}_{L^{(3,2)}} \underbrace{\begin{bmatrix} 19 \\ 1 \\ 25 \end{bmatrix}}_{b^{(2)}} = \underbrace{\begin{bmatrix} 19 \\ 1 \\ 21 \end{bmatrix}}_{b^{(3)}}.$$
$$b^{(3)} = (b_i^{(3)})_{i=1}^3$$

Metoda supstitucije

1. Iz treće jednadžbe je $x_3 = \frac{21}{7} = 3$.
2. Iz druge jednadžbe je $x_2 = \frac{1}{2}(1 + x_3) = 2$.
3. Iz prve jednadžbe je $x_1 = \frac{1}{5}(19 - x_2 - 4x_3) = 1$.

Opći prikaz metode

$$\left[\begin{array}{cccc|c} a_{11}^{(1)} & a_{12}^{(1)} & \cdots & a_{1n}^{(1)} & b_1^{(1)} \\ a_{21}^{(1)} & a_{22}^{(1)} & \cdots & a_{2n}^{(1)} & b_2^{(1)} \\ \vdots & \vdots & \cdots & \vdots & \vdots \\ a_{n1}^{(1)} & a_{n2}^{(1)} & \cdots & a_{nn}^{(1)} & b_n^{(1)} \end{array} \right]. \quad m_{i1} = \frac{a_{i1}^{(1)}}{a_{11}^{(1)}}, \quad i = 2, \dots, n.$$

$$\left[\begin{array}{cccc|c} a_{11}^{(1)} & a_{12}^{(1)} & \cdots & a_{1n}^{(1)} & b_1^{(1)} \\ 0 & a_{22}^{(2)} & \cdots & a_{2n}^{(2)} & b_2^{(2)} \\ \vdots & \vdots & \cdots & \vdots & \vdots \\ 0 & a_{n2}^{(2)} & \cdots & a_{nn}^{(2)} & b_n^{(2)} \end{array} \right].$$

Opći prikaz metode

$$m_{i2} = \frac{a_{i2}^{(2)}}{a_{22}^{(2)}}, \quad i = 3, \dots, n,$$

$$\left[\begin{array}{ccccc|c} a_{11}^{(1)} & a_{12}^{(1)} & \cdots & a_{1n}^{(1)} & | & b_1^{(1)} \\ a_{22}^{(2)} & \cdots & a_{2n}^{(2)} & | & b_2^{(2)} \\ \ddots & & \vdots & | & \vdots \\ & & a_{nn}^{(n)} & | & b_n^{(n)} \end{array} \right].$$

Uz prepostavku da je $a_{nn}^{(n)} \neq 0$, ovaj se linearni sustav lako rješava povratnom supstitucijom

pivotiranje

$$x_n = \frac{b_n^{(n)}}{a_{nn}^{(n)}},$$

$$x_i = \frac{1}{a_{ii}^{(i)}} \left(b_i^{(i)} - \sum_{j=i+1}^n a_{ij}^{(i)} x_j \right), \quad i = n-1, \dots, 1.$$

Uobičajeno **parcijalno pivotiranje** kao pivotni element bira element koji je po absolutnoj vrijednosti najveći u ostatku tog stupca — na glavnoj dijagonali ili ispod nje. Drugim riječima, ako je u k -tom koraku

$$|a_{rk}^{(k)}| = \max_{k \leq i \leq n} |a_{ik}^{(k)}|,$$

onda ćemo zamijeniti r -ti i k -ti redak i početi korak eliminacije elemenata k -tog stupca.

pivotiranje

Osim parcijalnog pivotiranja, može se provoditi i **potpuno pivotiranje**. U k -tom koraku, bira se maksimalni element u cijelom “ostatku” matrice $A^{(k)}$, a ne samo u k -tom stupcu. Ako je u k -tom koraku

$$|a_{rs}^{(k)}| = \max_{k \leq i, j \leq n} |a_{ij}^{(k)}|,$$

parcijalno pivotiranje - zamjena redaka

potpuno pivotiranje - zamjena redaka i stupaca

algoritam

Algoritam 5.2.1. (Gaussove eliminacije s parcijalnim pivotiranjem)

```
{Trokutasta redukcija}
for k := 1 to n - 1 do
    begin
        {Nađi maksimalni element u ostatku stupca}
        max_elt := 0.0;
        ind_max := k;
        for i := k to n do
            if abs(A[i, k]) > max_elt then
                begin
                    max_elt := abs(A[i, k]);
```

algoritam

```
ind_max := i;  
end;  
if max_elt > 0.0 then  
begin  
if ind_max <> k then  
{Zamijeni k-ti i ind_max-ti rijec  
begin  
for j := k to n do  
begin  
temp := A[ind_max, j];  
A[ind_max, j] := A[k, j];  
A[k, j] := temp;  
end;  
temp := b[ind_max];  
b[ind_max] := b[k];  
b[k] := temp;  
end;
```

```
for i := k + 1 to n do  
begin  
mult := A[i, k]/A[k, k];  
A[i, k] := 0.0; {Ne treba, ne koristi se kasnije}  
for j := k + 1 to n do  
A[i, j] := A[i, j] - mult * A[k, j];  
b[i] := b[i] - mult * b[k];  
end;  
end  
else  
{Matrica je singularna, stani s algoritmom}  
begin  
error := true;  
exit;  
end;  
end:
```

algoritam

```
{Povratna supstitucija, rješenje  $x$  ostavi u  $b$ }  
 $b[n] := b[n]/A[n, n];$   
for  $i := n - 1$  downto 1 do  
    begin  
         $sum := b[i];$   
        for  $j := i + 1$  to  $n$  do  
             $sum := sum - A[i, j] * b[j];$   
         $b[i] := sum/A[i, i];$   
    end;  
     $error := false;$ 
```

Zadatak 5.2.1. Pokušajte samostalno napisati algoritam koji koristi potpuno pivotiranje. Posebnu pažnju obratite na efikasno pamćenje zamjena varijabli koje su posljedica zamjena stupaca. Može li se isti princip efikasno primijeniti i za pamćenje zamjena redaka, tako da se potpuno izbjegnu eksplicitne zamjene elemenata u matrici A i vektoru b ?

gaussj.c

```
#include <math.h>
#define NRANSI
#include "nrutil.h"
#define SWAP(a,b) {temp=(a);(a)=(b);(b)=temp;}
void gaussj(float **a, int n, float **b, int m)
{
    int *indxc,*indxr,*ipiv;
    int i,icol,irow,j,k,l,ll;
    float big,dum,pivinv,temp;

    indxc=ivektor(1,n);
    indxr=ivektor(1,n);
    ipiv=ivektor(1,n);
    for (j=1;j<=n;j++) ipiv[j]=0;
    for (i=1;i<=n;i++) {
        big=0.0;
        for (j=1;j<=n;j++)
            if (ipiv[j] != 1)
                for (k=1;k<=n;k++) {
                    if (ipiv[k] == 0) {
                        if (fabs(a[j][k]) >= big) {
                            big=fabs(a[j][k]);
                            irow=j;
                            icol=k;
                        }
                    }
                }
        ++(ipiv[icol]);
        if (irow != icol) {
            for (l=1;l<=n;l++) SWAP(a[irow][l],a[icol][l])
            for (l=1;l<=m;l++) SWAP(b[irow][l],b[icol][l])
        }
    }

    indxr[i]=irow;
    indxc[i]=icol;
    if (a[icol][icol] == 0.0) nrerror("gaussj: Singular Matrix");
    pivinv=1.0/a[icol][icol];
    a[icol][icol]=1.0;
    for (l=1;l<=n;l++) a[icol][l] *= pivinv;
    for (l=1;l<=m;l++) b[icol][l] *= pivinv;
    for (ll=1;ll<=n;ll++)
        if (ll != icol) {
            dum=a[ll][icol];
            a[ll][icol]=0.0;
            for (l=1;l<=n;l++) a[ll][l] -= a[icol][l]*dum;
            for (l=1;l<=m;l++) b[ll][l] -= b[icol][l]*dum;
        }
    for (l=n;l>=1;l--) {
        if (indxr[l] != indxc[l])
            for (k=1;k<=n;k++)
                SWAP(a[k][indxr[l]],a[k][indxc[l]]);
    }
    free_ivektor(ipiv,1,n);
    free_ivektor(indxr,1,n);
    free_ivektor(indxc,1,n);
}

#undef SWAP
#undef NRANSI
```

LU faktorizacija

$$A^{(3)} = L^{(3,2)} L^{(3,1)} L^{(2,1)} A.$$

Matrica $A^{(3)}$ je gornjetrokutasta, a produkt $L^{(3,2)} L^{(3,1)} L^{(2,1)}$ je donjetrokutasta matrica. Dakle, polaznu matricu A smo množenjem slijeva donjetrokutastom matricom načinili gornjetrokutastom. To možemo pročitati i ovako:

$$A = LA^{(3)}, \quad L = (L^{(2,1)})^{-1} (L^{(3,1)})^{-1} (L^{(3,2)})^{-1},$$

gdje je L donjetrokutasta matrica. Lako se provjerava da je

$$L = \underbrace{\begin{bmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ 0 & 0 & 1 \end{bmatrix}}_{(L^{(2,1)})^{-1}} \underbrace{\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ -3 & 0 & 1 \end{bmatrix}}_{(L^{(3,1)})^{-1}} \underbrace{\begin{bmatrix} 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 4 & 1 \end{bmatrix}}_{(L^{(3,2)})^{-1}} = \begin{bmatrix} 1 & 0 & 0 \\ 2 & 1 & 0 \\ -3 & 4 & 1 \end{bmatrix}.$$

LU faktorizacija

Matrica A je produkt gornjotrokutaste i donjotrokutaste matrice

$$A = LA^{(3)} \quad U = A^{(3)} \quad A = LU.$$

Imamo: $A x = b$, $A = L U$, $L(U x) = b$, $U x = y$,

$$L y = b$$

$$U x = y$$

Linearni sustav riješen je pomoću 3 koraka:

1. Matricu sustava A treba faktorizirati u obliku $A = LU$, gdje je L donjetrokutasta, a U gornjetrokutasta matrica.
2. Rješavanjem donjetrokutastog sustava $Ly = b$ treba odrediti vektor $y = L^{-1}b$.
3. Rješavanjem gornjetrokutastog sustava $Ux = y$ treba odrediti vektor $x = U^{-1}y = U^{-1}(L^{-1}b)$.

supstitucija unaprijed

$$\begin{bmatrix} \ell_{11} & 0 & 0 & 0 \\ \ell_{21} & \ell_{22} & 0 & 0 \\ \ell_{31} & \ell_{32} & \ell_{33} & 0 \\ \ell_{41} & \ell_{42} & \ell_{43} & \ell_{44} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ b_3 \\ b_4 \end{bmatrix}.$$

linearni sustav s donjotrokutastom
matricom

matrica je regularna $\ell_{ii} \neq 0$

algoritam:

$$x_1 = \frac{b_1}{\ell_{11}};$$

za $i = 2, \dots, n$ {

$$x_i = \left(b_i - \sum_{j=1}^{i-1} \ell_{ij} x_j \right) / \ell_{ii}; \}$$

$$x_1 = \frac{b_1}{\ell_{11}}$$

$$x_2 = \frac{1}{\ell_{22}} (b_2 - \ell_{21} x_1)$$

$$x_3 = \frac{1}{\ell_{33}} (b_3 - \ell_{31} x_1 - \ell_{32} x_2)$$

$$x_4 = \frac{1}{\ell_{44}} (b_4 - \ell_{41} x_1 - \ell_{42} x_2 - \ell_{43} x_3).$$

supstitucija unazad

$$\begin{bmatrix} u_{11} & u_{12} & u_{13} & u_{14} \\ 0 & u_{22} & u_{23} & u_{24} \\ 0 & 0 & u_{33} & u_{34} \\ 0 & 0 & 0 & u_{44} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ b_3 \\ b_4 \end{bmatrix}$$

linearni sustav s gornjotrokutastom
matricom

algoritam

$$x_n = \frac{b_n}{u_{nn}};$$

za $i = n - 1, \dots, 1$ {

$$x_i = \left(b_i - \sum_{j=i+1}^n u_{ij}x_j \right) \Big/ u_{ii}; \}$$

$$x_4 = \frac{b_4}{u_{44}}$$

$$x_3 = \frac{1}{u_{33}} (b_3 - u_{34}x_4)$$

$$x_2 = \frac{1}{u_{22}} (b_2 - u_{23}x_3 - u_{24}x_4)$$

$$x_1 = \frac{1}{u_{11}} (b_1 - u_{12}x_2 - u_{13}x_3 - u_{14}x_4).$$

gauss eliminacija

SUBROUTINE gaussj(a,n,np,b,m,mp)

Linear equation solution by Gauss-Jordan elimination, equation (2.1.1) above.

a(1:n,1:n) is an input matrix stored in an array of physical dimensions **np by np**.

b(1:n,1:m) is an input matrix containing the m right-hand side vectors, stored in an array of physical dimensions **np by mp**. On output, **a(1:n,1:n)** is replaced by its matrix inverse, and **b(1:n,1:m)** is replaced by the corresponding set of solution vectors.

6 argumenta

void gaussj(float **a, int n, float **b, int m)

4 argumenta, ne treba NP i MP dimenzije od matrice i vektora.

b argument kod C verzije je matrica.

NR primjer

- ◆ Fortran
 - ◆ Files: Xgaussj.for, gaussj.for, MATRX1.dat
 - Xgaussj.for je modificirani "kod" originala xgaussj.for, problem je bio s čitanjem podataka iz file-a. xgaussj.for nemodificirani kod.
 - f77 -o fgauss Xgaussj.for gaussj.for
- ◆ C
 - ◆ Files: xgaussj.c gaussj.c nrutils/nrutil.c nrutils/, MATRX1.dat
 - ◆ nrutils/ sadrži .h fileove, tj. nrutil.h i nr.h
 - ◆ nrutil/, tj. nrutil.c i nrutil.h dijelovi za alociranje matrice, vektora itd..
 - gcc -o xgauss xgaussj.c gaussj.c nrutils/nrutil.c -I nrutils/

F77 gauss

fortran primjer: tgauss.f

Kompajliranje: g77 -o tg1 tgauss.f gaussj.for lib1.f

```
open (unit=8,file="linsys1.dat")
call readmatrix(MA1,n1,m1,fp)
call readvektor(b,L1,fp)
call gaussj(TA1,n1,NP,x,1,MP)
```

```
print *, "Solution is: "
call displayvektor(x,L1)
write(*,*) 'Inverse of Matrix A : '
call displaymatrix(TA1,n1,m1)
call multMatrixMatrix(TA1,MA1,XMM,n1,m1,n1,m1)
write(*,*) 'Inverse of Matrix A * A: '
call displaymatrix(XMM,n1,m1)
```

C gauss

C primjer: tgauss.c

Kompajliranje: gcc -g -o tg tgauss.c libnr.c gaussj.c nrutils/nrutil.c -I nrutils/
ako se nrutil.c i header fileovi nalaze u poddirektoriju nrutils

```
#include "nr.h"
#include "nrutil.h"
#include "libnr.h"
#define NP 20
-----
float *b,*x;
/* float b[10],x[10]; */
float **MA1,**TA1,**XMM,**X;
```

```
X=matrix(1,NP,1,NP);
b=vector(1,NP);
readmatrix(MA1,n1,m1,fp);
readvektor(b,L1,fp);

gaussj(TA1,n1,X,1);
printf("Solution is:\n");
displaymatrix(X,n1,1);
printf("Inverse of A :\n"); displaymatrix(TA1,n1,m1);
```

LU metoda

SUBROUTINE ludcmp(a,n,np,indx,d)

void ludcmp(float **a, int n, int *indx, float *d)

Given a matrix **a[1..n][1..n]**, this routine replaces it by the LU decomposition of a rowwise permutation of itself. **a** and **n** are input. **a** is output; **indx[1..n]** is an output vector that records the row permutation effected by the partial pivoting; **d** is output as ± 1 depending on whether the number of row interchanges was even or odd, respectively. This routine is used in combination with **lubksb** to solve linear equations or invert a matrix.

NR PRIMJER:

```
gcc -o xludcmpC xludcmp.c ludcmp.c nrutils/nrutil.c -I nrutils/
```

```
g77 -o xludcmpF xludcmp.for ludcmp.for
```

```
gcc -o xlubsub xlubksb.c lubksb.c ludcmp.c nrutils/nrutil.c -I nrutils/
```

```
g77 -o xlubsubF xlubksb.for lubksb.for ludcmp.for
```

SUBROUTINE lubksb(a,n,np,indx,b)

void lubksb(**float** **a, **int** n, **int** *indx, **float** b[])

Solves the set of n linear equations $A \cdot X = B$. Here **a** is input, not as the matrix A but rather as its LU decomposition, determined by the routine ludcmp. **indx** is input as the permutation vector returned by ludcmp. **b(1:n)** is input as the right-hand side vector B, and returns with the solution vector X. **a**, **n**, **np**, and **indx** are not modified by this routine and can be left in place for successive calls with different right-hand sides b. This routine takes into account the possibility that **b** will begin with many zero elements, so it is efficient for use in matrix inversion.

g77 -g -o tludecmp tludecmp.for ludcmp.for lib1.f

gcc -g -o cludecmp tludecmp.c ludcmp.c nrutils/nrutil.c libnr.c -I nrutils/

gcc -g -o cludecmp tludecmp.c ludcmp.c nrutils/nrutil.c libnr.c -I nrutils/

g77 -o tlubksub tlubksub.for lubksb.for ludcmp.for lib1.f

gcc -g -o clubksb tlubksub.c ludcmp.c lubksb.c nrutils/nrutil.c libnr.c -I nrutils/

Osobine metoda

- ❖ Kvadratični linearni sustavi, broj nepoznanica= jednadžbi
- ❖ Pivotiranje je neophodno za stabilnost algoritma
- ❖ Gaussova eliminacija
 - ❖ Slabosti
 - ❖ Zahtjeva desnu stranu (vektor b), koju treba mjenjati u algoritmu
 - ❖ Sporost, naročito kada nam ne treba inverzna matrica
 - ❖ Prednost, jednostavna
- ❖ LU faktorizacija, manji broj operacije od Gauss metode

Zadaci za praktikum

1. Izvršite programe napravljene od xgaussj i tgaussj rutina koje koriste gaussovou eliminaciju.
2. Napravite program koji pročita linearni sustav iz datoteke LIN.DAT i riješite sustav gaussovom metodom.
3. Izvršite programe dobivene iz xludcmp , xlubksb, tludecmp, tlubksb koji koriste LU faktorizaciju.
4. Napravite program koji pročita linearni sustav iz datoteke LIN.DAT i riješite sustav LU faktorizacijom.
5. Dobiveno rješenje pošaljite kao poruku u mailu (jedno rješenje), a source programa stavite kao attachment. Mail je
Aleksandar.Maksimovic@irb.hr

Literatura

- ◆ Online literatura:
 - ◆ Numerička matematika-osnovni udžbenik, PMF, projekt mzt.
 - ◆ Numerical Recipes in C
 - ◆ Numerical Recipes in Fortran